ROBUST BINARY FUSED COMPRESSIVE SENSING USING ADAPTIVE OUTLIER PURSUIT

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ABSTRACT

We propose a new method, robust binary fused compressive sensing (RoBFCS), to recover sparse piece-wise smooth signals from 1-bit compressive measurements. The proposed method is a modification of our previous binary fused compressive sensing (BFCS) algorithm, which is based on the binary iterative hard thresholding (BIHT) algorithm. As in BIHT, the data term of the objective function is a one-sided ℓ_1 (or ℓ_2) norm. Experiments show that the proposed algorithm is able to take advantage of the piece-wise smoothness of the original signal and detect sign flips and correct them, achieving more accurate recovery than BFCS and BIHT.

Index Terms— 1-bit compressive sensing, iterative hard thresholding, group sparsity, signal recovery.

1. INTRODUCTION

In *compressive sensing* (CS) [1], [2], a sparse signal $\mathbf{x} \in \mathbb{R}^n$ is shown to be recoverable from a few linear measurements

$$\mathbf{b} = \mathbf{A}\mathbf{x},\tag{1}$$

where $\mathbf{b} \in \mathbb{R}^m$ (with m < n), $\mathbf{A} \in \mathbb{R}^{m \times n}$ is the sensing matrix (which satisfies some conditions), and the fact that m < n makes (1) an ill-posed problem. This classical formulation assumes real-valued measurements, thus ignoring that, in reality, any acquisition involves quantization. In *quantized CS* (QCS) [3], [4], [5], [6], [7], [8], this fact is taken into account. An interesting extreme case of QCS is 1-bit CS [9],

$$\mathbf{y} = \operatorname{sign}(\mathbf{A}\mathbf{x}),\tag{2}$$

where $sign(\cdot)$ is the element-wise sign function. Such 1-bit measurements can be acquired by a comparator with zero, which is very inexpensive and fast, as well as robust to amplification distortions. In contrast with the measurement model of standard CS, 1-bit measurements are blind to the magnitude of the original signal \mathbf{x} ; the goal may then only be to recover \mathbf{x} , up to an intrinsically unrecoverable magnitude.

The first algorithm for signal recovery from 1-bit measurements, named renormalized fixed point iteration (RFPI)

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was proposed in [9]. Soon after, [10] showed that recovery from nonlinearly distorted measurements is also possible, and [10] proposed a greedy algorithm (matching sign pursuit) [11]. After that seminal work, several algorithms for 1-bit CS have appeared; a partial list includes linear programming [12], [13], restricted-step shrinkage [14], and binary iterative hard thresholding (BIHT), which performs better than the previous algorithms. Algorithms for 1-bit CS, based on generalized approximate message passing [15] and majorization-minimization [16], were proposed in [17] and [18], respectively. In [19], the ℓ_1 norm in the data-term of [12] was replaced by an ℓ_0 norm; the resulting problem is solved by successive approximations, yielding a sequence of simpler problems, not requiring prior knowledge about the sparsity of the original signal. Considering the possibility of sign flips due to by noise, [20] introduced the adaptive outlier pursuit (AOP) algorithm, and [21] extended it into an algorithm termed noise-adaptive RFPI, which doesn't need prior information on the signal sparsity and number of sign flips.

More recently, [22] and [23] applied 1-bit CS in image acquisition, [24] studied matrix completion from noisy 1-bit observations, [25] used methods of statistical mechanics to examine typical properties of 1-bit CS. The authors of [26] addressed 1-bit CS using their recent *gradient support pursuit* (GraSP) [27] algorithm; finally, a *quantized iterative hard thresholding* method proposed in [28] provides a bridge between 1-bit and high-resolution QCS.

Recently, we proposed *binary fused compressive sensing* (BFCS) [29], [30] to recover group-sparse signals from 1-bit CS measurements. The rationale is that group-sparsity may express more structured knowledge about the unknown signal than simple sparsity, thus potentially allowing for more robust recovery from fewer measurements. In this paper, we further consider the possibility of sign flips, and propose *robust BFCS* (RoBFCS) based on the AOP method [20].

The rest of the paper is organized as follows: Section II reviews the BIHT and BFCS algorithms, and introduces the proposed RoBFCS method; Section III reports experimental results and Section IV concludes the paper.

2. ROBUST BFCS

2.1. The Observation Model

In this paper, we consider the noisy 1-bit measurement model,

$$\mathbf{y} = \operatorname{sign} \left(\mathbf{A} \mathbf{x} + \mathbf{w} \right), \tag{3}$$

where $\mathbf{y} \in \{+1, -1\}^m$, $\mathbf{A} \in \mathbb{R}^{m \times n}$ is as above, $\mathbf{x} \in \mathbb{R}^n$ is the original signal, and $\mathbf{w} \in \mathbb{R}^m$ is additive white Gaussian noise with the variance σ^2 , due to which some of the measurements signs may change with the respect to the noiseless measurements as given by (2).

2.2. Binary Iterative Hard Thresholding (BIHT)

To recover x from y, Jacques et al [31] proposed the criterion

$$\min_{\mathbf{x}} f(\mathbf{y} \odot \mathbf{A}\mathbf{x})$$
subject to $\|\mathbf{x}\|_2 = 1, \ \mathbf{x} \in \Sigma_K$,

where " \odot " represents the Hadamard (element-wise) product, $\Sigma_K = \{\mathbf{x} \in \mathbb{R}^n : \|\mathbf{x}\|_0 \le K\}$ (with $\|\mathbf{v}\|_0$ denoting the number of non-zero components in \mathbf{v}) is the set of K-sparse signals, and f is one of the penalty functions defined next. To penalize linearly the sign consistency violations, the choice is $f(\mathbf{z}) = 2 \|\mathbf{z}_-\|_1$, where $\mathbf{z}_- = \min(\mathbf{z}, 0)$ (where the minimum is applied component-wise and the factor 2 is included for later convenience) and $\|\mathbf{v}\|_1 = \sum_i |v_i|$ is the ℓ_1 norm. Quadratic penalization of the sign violations is achieved by using $f(\mathbf{z}) = \frac{1}{2} \|\mathbf{z}_-\|_2^2$, where the factor 1/2 is also included for convenience. The *iterative hard thresholding* (IHT) [32] algorithm applied to (4) (ignoring the norm constraint during the iterations) leads to the BIHT algorithm [31]:

Algorithm BIHT

- 1. set $t = 0, \tau > 0, \mathbf{x}_0$ and K
- 2. repeat
- 3. $\mathbf{v}_{t+1} = \mathbf{x}_t \tau \partial f \left(\mathbf{y} \odot \mathbf{A} \mathbf{x}_t \right)$
- 4. $\mathbf{x}_{t+1} = \mathcal{P}_{\Sigma_K} \left(\mathbf{v}_{t+1} \right)$
- 5. $t \leftarrow t + 1$
- 6. **until** some stopping criterion is satisfied.
- 7. **return** $\mathbf{x}_t / \|\mathbf{x}_t\|$

In this algorithm, ∂f denotes the subgradient of the objective (see [31], for details), which is given by

$$\partial f\left(\mathbf{y} \odot \mathbf{A} \mathbf{x}\right) = \begin{cases} \mathbf{A}^{T} \left(\operatorname{sign}(\mathbf{A} \mathbf{x}) - \mathbf{y}\right), & \ell_{1} \text{ penalty} \\ \left(\mathbf{Y} \mathbf{A}\right)^{T} \left(\mathbf{Y} \mathbf{A} \mathbf{x}\right)_{-}, & \ell_{2} \text{ penalty}, \end{cases}$$
(5)

where $\mathbf{Y} = \operatorname{diag}(\mathbf{y})$ is a diagonal matrix with vector \mathbf{y} in its diagonal. Step 3 corresponds to a sub-gradient descent step (with step-size τ), while Step 4 performs the projection onto the non-convex set Σ_K , which corresponds to computing the best K-term approximation of \mathbf{v} , *i.e.*, keeping the K largest

components in magnitude and setting the others to zero. Finally, the returned solution is projected onto the unit sphere to satisfy the constraint $\|\mathbf{x}\|_2 = 1$. The versions of BIHT for the ℓ_1 and ℓ_2 penalties are referred to as BIHT and BIHT- ℓ_2 , respectively.

2.3. Binary Fused Compressive Sensing (BFCS)

We begin by introducing some notation. The TV semi-norm of a vector $\mathbf{v} \in \mathbb{R}^n$ is given by $\mathrm{TV}(\mathbf{v}) = \sum_{i=1}^{n-1} |v_{i+1} - v_i|$. For $\varepsilon \geq 0$, we denote as T_ε the ε -radius TV ball, i.e., $T_\varepsilon = \{\mathbf{v} \in \mathbb{R}^p : \mathrm{TV}(\mathbf{v}) \leq \varepsilon\}$. The projection onto T_ε (denoted $\mathcal{P}_{T_\varepsilon}$) can be computed by the algorithm proposed in [33]. Let $\mathfrak{G}(\mathbf{v}) = \bigcup_{i=1}^{\mathcal{K}(\mathbf{v})} \mathcal{G}_i(\mathbf{v})$, where each $\mathcal{G}_k(\mathbf{v}) \subset \{1,...,n\}$ is a set of consecutive indices $\mathcal{G}_k(\mathbf{v}) = \{i_k,...,i_k + |\mathcal{G}_k| - 1\}$ such that, for $j \in \mathcal{G}_k$, $v_j \neq 0$, while $v_{i_k-1} = 0$ and $v_{i_k+|\mathcal{G}_k|} = 0$ (assume that $v_0 = v_{n+1} = 0$); $\mathcal{G}_k(\mathbf{v})$ is the k-th group of indices of consecutive non-zero components of \mathbf{v} , and there are $\mathcal{K}(\mathbf{v})$ such groups. Let $\mathbf{v}_{\mathcal{G}_k} \in \mathbb{R}^{|\mathcal{G}_k|}$ be the sub-vector of \mathbf{v} with indices in \mathcal{G}_k .

Obviously, the criterion in (4) doesn't encourage groupsparsity. To achieve that goal, we propose the criterion

$$\min_{\mathbf{x}} f(\mathbf{y} \odot \mathbf{A} \mathbf{x})$$
subject to $\|\mathbf{x}\|_{2} = 1$, $\mathbf{x} \in \Sigma_{K} \cap S_{\epsilon}$

where S_{ϵ} is defined as

$$S_{\epsilon} = \left\{ \mathbf{x} \in \mathbb{R}^n : \overline{\text{TV}}(\mathbf{x}_{\mathcal{G}_k}) \le \epsilon, k = 1, \cdots, \mathcal{K}(\mathbf{x}) \right\}.$$
 (7)

where $\overline{\text{TV}}(\mathbf{x}_{\mathcal{G}_k}) = (|\mathcal{G}_k| - 1)^{-1} \, \text{TV}(\mathbf{x}_{\mathcal{G}_k})$ is a normalized TV, where $|\mathcal{G}_k| - 1$ is the number of absolute differences in TV $(\mathbf{x}_{\mathcal{G}_k})$. In contrast with a standard TV ball, S_ϵ promotes the "fusion" of components only inside each non-zero group, that is, the TV regularizer does not "compete" with the sparsity constraint imposed by $\mathbf{x} \in \Sigma_K$.

To address the optimization problem in (6), we propose the following algorithm (which is a modification of BIHT):

Algorithm BFCS

- 1. set $t = 0, \tau > 0, \epsilon > 0, K$ and \mathbf{x}_0
- 2. repeat
- 3. $\mathbf{v}_{t+1} = \mathbf{x}_t \tau \partial f \left(\mathbf{y} \odot \mathbf{A} \mathbf{x}_t \right)$
- 4. $\mathbf{x}_{t+1} = \mathcal{P}_{S_{\epsilon}}(\mathcal{P}_{\Sigma_K}(\mathbf{v}_{t+1}))$
- 5. $t \leftarrow t + 1$
- 6. **until** some stopping criterion is satisfied.
- 7. return $\mathbf{x}_t / \|\mathbf{x}_t\|$

Notice that the composition of projections in line 4 is not in general equal to the projection on the non-convex set $\Sigma_K \cap S_{\epsilon}$, *i.e.*, $\mathcal{P}_{\Sigma_K \cap S_{\epsilon}} \neq \mathcal{P}_{S_{\epsilon}} \circ \mathcal{P}_{\Sigma_K}$. However, this composition does satisfy some relevant properties, which result from the structure of $\mathcal{P}_{S_{\epsilon}}$ expressed in the following lemma (the proof of which is quite simple, but is omitted due to lack of space).

Lemma 1 Let $\mathbf{v} \in \mathbb{R}^n$ and $\mathbf{x} = \mathcal{P}_{S_{\epsilon}}(\mathbf{v})$, then

$$\mathbf{x}_{\mathcal{G}_{k}} = \mathcal{P}_{T_{\epsilon(|\mathcal{G}_{k}|-1)}}(\mathbf{v}_{\mathcal{G}_{k}}), \quad \text{for } k = 1, \cdots, \mathcal{K}(\mathbf{v});$$

$$\mathbf{x}_{\overline{\mathfrak{G}}(\mathbf{v})} = \mathbf{0}, \tag{8}$$

where $\overline{\mathfrak{G}}(\mathbf{v}) = \{1, \dots, n\} \setminus \mathfrak{G}(\mathbf{v})$ and $\mathbf{0}$ is a vector of zeros.

The other relevant property of $\mathcal{P}_{S_{\epsilon}}$ is that it preserves sparsity, as expressed formally in the following lemma.

Lemma 2 If $\mathbf{v} \in \Sigma_K$, then $\mathcal{P}_{S_{\epsilon}}(\mathbf{v}) \in \Sigma_K$. Consequently, for any $\mathbf{v} \in \mathbb{R}^n$, $\mathcal{P}_{S_{\epsilon}}(\mathcal{P}_{\Sigma_K}(\mathbf{v})) \in \Sigma_K \cap S_{\epsilon}$.

That is, although it is **not** guaranteed that $\mathcal{P}_{S_{\epsilon}}(\mathcal{P}_{\Sigma_K}(\mathbf{v}))$ coincides with the orthogonal projection of \mathbf{v} onto $\Sigma_K \cap S_{\epsilon}$, it belongs to this non-convex set. In fact, the projection onto $\Sigma_K \cap S_{\epsilon}$ can be shown to be an NP-hard problem [34], since it belongs to the class of shaped partition problems [35], [36] with variable number of parts.

2.4. Proposed Formulation and Algorithm

In this paper we extend the BFCS approach to deal with the case where there may exist some sign flips in the measurements. To this end, we adopt the AOP technique [20], yielding a new approach that we call robust BFCS (RoBFCS); the similarly robust version of BIHT is termed RoBIHT. Assume that there are at most L sign flips and define the binary vector $\Lambda \in \{-1, +1\}^m$ as

$$\Lambda_i = \begin{cases}
-1 & \text{if } y_i \text{ is "flipped"}; \\
+1 & \text{otherwise.}
\end{cases}$$
(9)

Then, the criterion of RoBFCS is given by

$$\min_{\mathbf{x} \in \mathbb{R}^{n}, \ \mathbf{\Lambda} \in \{-1, +1\}^{m}} f\left(\mathbf{y} \odot \mathbf{\Lambda} \odot \mathbf{A} \mathbf{x}\right)$$
subject to $\|\mathbf{x}\|_{2} = 1, \ \mathbf{x} \in \Sigma_{K} \cap S_{\epsilon}$

$$\|\mathbf{\Lambda}_{-}\|_{1} \leq L,$$
(10)

where $\Lambda_{-} = \min{\{\Lambda, 0\}}$. Problem (10) is mixed continuous/discrete, and clearly difficult. A natural approach to address (10) is via alternating minimization, as follows.

Algorithm Framework of RoBFCS

- set $t = 0, \Lambda_0 = \mathbf{1} \in \mathbb{R}^m, \epsilon > 0, K, L$ and \mathbf{x}_0 1.
- 2.
- 3. $\mathbf{x}_{t+1} = \Phi\left(\mathbf{y} \odot \mathbf{\Lambda}_t, K, \epsilon\right)$
- $\mathbf{\Lambda}_{t+1} = \Psi \left(\mathbf{y} \odot \mathbf{A} \mathbf{x}_{t+1}, L \right)$ 4.
- 5. $t \leftarrow t + 1$
- until some stopping criterion is satisfied. 6.
- return $\mathbf{x}_t / \|\mathbf{x}_t\|$

In this algorithm (template) lines 3 and 4 correspond to minimizing (10) with respect to x and Λ , respectively. The minimization w.r.t. x defines the function

$$\Phi\left(\mathbf{u}, K, \epsilon\right) = \underset{\mathbf{x} \in \mathbb{R}^{n}}{\min} \ f\left(\mathbf{u} \odot \mathbf{A} \mathbf{x}\right)$$
subject to $\|\mathbf{x}\|_{2} = 1, \ \mathbf{x} \in \Sigma_{K} \cap S_{\epsilon}$

which is an instance of (6). The minimization w.r.t. Λ defines the function

$$\Psi\left(\mathbf{z}, L\right) = \arg\min_{\mathbf{\Lambda} \in \{-1, 1\}^{m}} f\left(\mathbf{z} \odot \mathbf{\Lambda}\right)$$
subject to $\|\mathbf{\Lambda}_{-}\|_{1} \leq L$ (12)

As shown in [20], [21], function (12) is given in closed form by

$$\left(\Psi\left(\mathbf{z},L\right)\right)_{i} = \begin{cases} -1 & \text{if } z_{i} \geq \tau; \\ +1 & \text{otherwise,} \end{cases}$$
 (13)

where τ is the L-th largest element (in magnitude) of z.

In the proposed RoBFCS algorithm, rather than implementing (11) by running the BFCS algorithm until some stopping criterion is satisfied, a single step thereof is applied, followed by the implementation of (12) given by (13).

Algorithm RoBFCS

- 1. set $t = 0, \tau > 0, \epsilon > 0, K, L$ and $\mathbf{x}_0, \mathbf{\Lambda}_0 = \mathbf{1} \in \mathbb{R}^m$
- 2.
- $\mathbf{v}_{t+1} = \mathbf{x}_k \tau \partial f \left(\mathbf{y} \odot \mathbf{\Lambda}_t \odot \mathbf{A} \mathbf{x}_k \right)$ 3.
- 4.
- $\mathbf{x}_{t+1} = \mathcal{P}_{S_{\epsilon}} \left(\mathcal{P}_{\Sigma_{K}} (\mathbf{v}_{t+1}) \right)$ $\mathbf{\Lambda}_{t+1} = \Psi \left(\mathbf{y} \odot \mathbf{A} \mathbf{x}_{t+1}, L \right)$ 5.
- 6.
- 7. until some stopping criterion is satisfied.
- return $\mathbf{x}_t / \|\mathbf{x}_t\|$

The subgradient in line 3 is as given by (5), with y replaced with $\mathbf{y} \odot \mathbf{\Lambda}_t$. If the original signal is known to be nonnegative, the algorithm includes a projection onto \mathbb{R}^n_+ in each iteration, *i.e.*, line 4 becomes $\mathbf{x}_{k+1} = \mathcal{P}_{\mathbb{R}^n_+} (\mathcal{P}_{S_{\epsilon}} (\mathcal{P}_{\Sigma_K} (\mathbf{v}_{k+1})))$. The versions of RoBFCS (RoBIHT) with ℓ_1 and ℓ_2 objectives are referred to as RoBFCS and RoBFCS-\$\ell_2\$ (RoBIHT and RoBIHT- ℓ_2), respectively.

3. EXPERIMENTS

In this section, we report results of experiments aimed at studying the performance of RoBFCS. All the experiments were performed using MATLAB on a 64-bit Windows 7 PC with an Intel Core i7 3.07 GHz processor. In order to measure the performance of different algorithms, we employ the following five metrics defined on an estimate e of an original vector **x** (both of unit norm):

- Mean absolute error, $\mathbf{MAE} = \|\mathbf{x} \mathbf{e}\|_1 / n$;
- Mean square error, $\mathbf{MSE} = \|\mathbf{x} \mathbf{e}\|_2^2 / n$;
- Position error rate, $\mathbf{PER} = \sum_{i} ||\operatorname{sign}(x_i)| |\operatorname{sign}(e_i)|| / n;$
- Angle error, $\mathbf{AE} = \arccos \langle \mathbf{x}, \mathbf{e} \rangle / \pi$;

BIHT RoBIHT Metrics BIHT- ℓ_2 **BFCS** BFCS- ℓ_2 RoBIHT- ℓ_2 RoBFCS RoBFCS- ℓ_2 0.0019 0.0001 MAE 0.0019 0.0032 0.0008 0.0034 0.0038 0.0038 MSE 7.43E-5 1.65E-4 2.87E-5 1.78E-4 7.12E-5 2.04E-4 4.00E-7 2.06E-4 PER 1.8% 0.9% 4.9% 2.0% 4.7% 0% 5.2% 4.1%HE 0.0450 0.1360 0.0530 0.0995 0.0050 0.1420 0.0010 0.1390 ΑE 0.1234 0.1852 0.0764 0.1927 0.1208 0.2070 0.0085 0.2082

Table 1. Experimental results

• Hamming error, $\mathbf{HE} = \|\operatorname{sign}(\mathbf{Ax}) - \operatorname{sign}(\mathbf{Ae})\|_0/m$. The original signals \mathbf{x} are taken as sparse and piece-wis

The original signals ${\bf x}$ are taken as sparse and piece-wise smooth, of length n=2000 with sparsity level K=160; specifically,

$$\bar{x}_{i} = \begin{cases} 10 + 0.1 \, k_{i}, & i \in \bigcup_{i=1}^{0.25d} \mathcal{B}_{i} \\ 15 + 0.1 \, k_{i}, & i \in \bigcup_{i=0.25d+1}^{0.5d} \mathcal{B}_{i} \\ -10 + 0.1 \, k_{i}, & i \in \bigcup_{i=0.5d+1}^{0.75d} \mathcal{B}_{i} \\ -15 + 0.1 \, k_{i}, & i \in \bigcup_{i=0.75d+1}^{d} \mathcal{B}_{i} \\ 0, & i \notin \bigcup_{i=1}^{d} \mathcal{B}_{i} \end{cases}$$
(14)

where the k_i are independent samples of a zero-mean, unit variance Gaussian random variable, d is the number of nonzero groups of \mathbf{x} , and $\mathcal{B}_i, i \in \{1, \dots, d\}$ indexes the i-th group, defined as

$$\mathcal{B}_i = \{50 + (i-1) n/d + 1, \dots, 50 + (i-1) n/d + K/d\}.$$

The signal is then normalized, $\mathbf{x} = \bar{\mathbf{x}}/\|\bar{\mathbf{x}}\|_2$. The sensing matrix \mathbf{A} is a 2000×2000 matrix with components sampled from the standard normal distribution. Finally, observations \mathbf{y} are obtained by (3), with noise standard deviation $\sigma=1$. The assumed number of sign flips is L=10.

We run the algorithms BIHT, BIHT- ℓ_2 , BFCS, BFCS- ℓ_2 , RoBIHT, RoBIHT- ℓ_2 , RoBFCS and RoBFCS- ℓ_2 . The stopping criterion is $\|\mathbf{x}_{(k+1)} - \mathbf{x}_{(k)}\| / \|\mathbf{x}_{(k+1)}\| \le 0.001$, where $\mathbf{x}_{(k)}$ is estimate at the k-th iteration. Following the setup of [31] and [20], the step-size of BIHT and RoBIHT and that of BIHT- ℓ_2 and Ro BIHT- ℓ_2 is $\tau=1$ and 1/m, respectively. While in BFCS, BFCS- ℓ_2 , RoBFCS, RoBFCS- ℓ_2 , τ and ϵ are hand tuned for the best improvement in SNR. The quantitative results are shown in Table 1.

From Table 1, we can see that RoBFCS performs the best in terms of the metrics considered. Moreover, the algorithms with ℓ_1 barrier perform better than those with ℓ_2 barrier.

4. CONCLUSIONS

Based on the previously proposed BFCS (binary fused compressive sensing) method, we have proposed an algorithm for recovering sparse piece-wise smooth signals from 1-bit compressive measurements with some sign flips. We have shown that if the original signals are in fact sparse and piece-wise smooth and there are some sign flips in the measurements, the

proposed method (termed RoBFCS – robust BFCS) outperforms (under several accuracy measures) the previous methods BFCS and BIHT (binary iterative hard thresholding). Future work will aim at making RoBFCS adaptive in terms of K and L.

5. REFERENCES

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